Kolmogorov Complexity

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Joseph Biberstine

How complex are these messages

- Examples
 - \odot 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20
 - $_{\odot}$ 1 3 5 7 9 11 13 15 17 19 21 23 25 27 29 31 33 35 37 39

\circ 1 2 1 4 1 2 1 8 1 2 1 4 1 2 1 16 1 2 1 4 1 2 1 8 1 2 1 4 1 2

- Do they seem increasingly complex?
 - \circ How can we quantify this?
 - "Complexity" is a word with many meanings..
 - but one thought to pursue might be:
 - They seem increasingly difficult to describe

Definition

- Kolmogorov complexity = descriptive complexity
- First, we need to pick a language strong enough to describe all messages of interest
 - Turing-completeness allows all computable messages
 - Let's use English as an accessible example language
 - In many cases, we will find descriptions shorter than simply restating the message
- The Kolmogorov complexity of a message is defined to be the length of its shortest description in this language
 Description should be exact and unambiguous

Kolmogorov complexity of examples

• 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 • integers from 1 to 20

Kolmogorov complexity = Length = 21

1 3 5 7 9 11 13 15 17 19 21 23 25 27 29 31 33 35 37 39
odd integers from 1 to 39
Length = 25

• 1 2 1 4 1 2 1 8 1 2 1 4 1 2 1 16 1 2 1 4 1 2 1 8 1 2 1 4 1 2

 write "1 ", then repeatedly concatenate your work to itself and double the last number. take first 30 of result.

■ Length = 113

\circ greatest power of 2 dividing integers from 1 to \Box 30

■ Length = 50

Puzzles

How complex (in our sense) are these?
 0 1 4 0 8 4 5 0 7 0 4 2 2 5 3 5 2 1 1 2 6 7 6 0 5 6 3 3

Solutions(?)

- $\begin{array}{c} \bullet \ 0 \ 1 \ 0 \ 3 \ 0 \ 9 \ 2 \ 7 \ 8 \ 3 \ 5 \ 0 \ 5 \ 1 \ 5 \ 4 \ 6 \ 3 \ 9 \ 1 \ 7 \ 5 \ 2 \ 5 \ 7 \ 7 \ 3 \ 1 \ 9 \ 5 \ 8 \ 7 \\ 6 \ 2 \ 8 \ 8 \ 6 \ 5 \ 9 \ 7 \ 9 \ 3 \ 8 \ 1 \ 4 \ 4 \ 3 \ 2 \ 9 \ 8 \ 9 \ 6 \ 9 \ 0 \ 7 \ 2 \ 1 \ 6 \ 4 \ 9 \ 4 \ 8 \ 4 \ 5 \end{array}$
 - o 1st 64 digits of 1/97
 - But maybe we could do better?
 - Could only be certain if we checked every smaller description
- 12548906733029750954218637379654
 71023815072396844180793619540872
 ?
 - \circ 64 random digits
 - \circ Is there a shorter description than the sequence itself?
 - Who knows?

Languages: Natural

- Any Turing-complete language could serve as the basis
- But which one?
- English is convenient, but it's not rigorous for this purpose
 o the number of people on Earth January 1st 2000 at midnight GMT
 - Does this describe exactly one numeric message?
 - What about people jumping or flying ...
 - Being born or dying or ...

Languages: Formal

- Formal languages solve this problem, but they are less intuitive than natural language
 - How do we find descriptions shorter than a literal restatement?
 - And if we do, how can we be certain that any one we've found is the shortest?

In practice

- Though Kolmogorov complexity is tricky to calculate exactly, except perhaps for the shortest of strings, we can get reasonable approximations
 - Consider a compression algorithm such as one used to create zip archives
 - It will produce descriptions of any input, trying very thoroughly to be as succinct as possible
 - Such algorithms are as good an attempt as we are likely to find in full generality

What does it measure?

- Now that we have this concept defined well and on a solid theoretical footing using a T-C formal language.. what's it telling us, anyway?
- We call it "complexity", but can we say more specifically what it means? Following Feldman, it:
 - measures "randomness"
 - incompressibility/unstructuredness
 - o does not measure "pattern or structure or correlation or organization"

In artificial life

- Could we use Kolmogorov complexity to guide evolution in a simulation?
 - Selection could be wholly or partially determined by the complexity of the genome
 - Minimize it: Maybe we'll get elegantly simple agents!
 Instead: Boringness
 - Maximize it: Maybe we'll get complex agents!
 - Instead: Randomness
 - Instead of generating agents with behavior that is complex in another, more desirable, sense of the word, we would more likely find a structureless noise

So it's a wash?

- Perhaps not; consider these experiments involving a fitness function:
 - When it begins to plateau across the population, add a component favoring simpler genomes
 - Maybe we can drive inessential complexity out of the solution
 - Add a small component favoring complexity and reduce the global mutation rate
 - Instead of changing alleles blindly, maybe we will occasionally pinpoint those areas that will effect the most ambitious search
- What do you think would happen?

Interesting properties

- How great can a message's Kolmogorov complexity be?

 Always bounded by whatever literal restatement we could lazily give it, potentially plus a little padding amounting to "the following is literal:"
 - This is important, if we ignored it and encoded the message integers between 1 and 50 as integers between 1 and 50 we will probably cause confusion

Confusing descriptions with literal messages

E-MAIL ADDRESSES IT WOULD BE REALLY ANNOYING TO GIVE OUT OVER THE PHONE.

BY MICHAEL WARD

MikeUnderscore2004@yahoo.com

MikeAtYahooDotCom@hotmail.com

Mike_WardAllOneWord@yahoo.com

AAAAAThatsSixAs@yahoo.com

One1TheFirstJustTheNumberTheSecondSpelledOut@hotmail.com

http://www.mcsweeneys.net/links/lists/27MichaelWard.html

Interesting properties

- For theoretical purposes, the base language used doesn't matter too much
 - For any two potential base languages, there is a constant upper bound on the difference between the complexities calculated using either
 - Constant never depends on a particular message, only on the languages
 - Still, the constant almost surely enormous

Interesting properties

- Joint Kolmogorov complexity satisfies an equality reminiscent of joint entropy
 - $\circ H(X,Y) = H(X) + H(Y \mid X)$
 - $\circ K(X,Y) = K(X) + K(Y \mid X) + O(\log(K(X,Y)))$
 - \circ Joint means we are to describe both
 - Conditional means we can use the given message as input
 - Consider X = **01** and Y = **0101**
 - Our description for calculating K(Y | X) could be X twice
 - Our description for calculating K(X, Y) could be 01, then the first two characters twice
 - Logarithmic factor essentially allows for space in the description to explain which part encodes message X and which encodes message Y

That's all

